

SCert: Speculative Certification in Replicated Software Transactional Memories

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Roadmap

Motivation

Related Work

SCert

Examples

Results

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Transactional Memory

- Set of mechanisms for shared memory access
- Uses de concept of *Transaction*

Programmers only indicate the set of operations that must be performed atomically: simpler than using Locks explicitly

Distributed Transactional Memory

Provides fault tolerance and increased performance

DSTM vs Distributed Shared Memory

- *Similar*: Hides the distribution from the programmers
- *Different*: Synchronization is only performed at the transaction boundaries

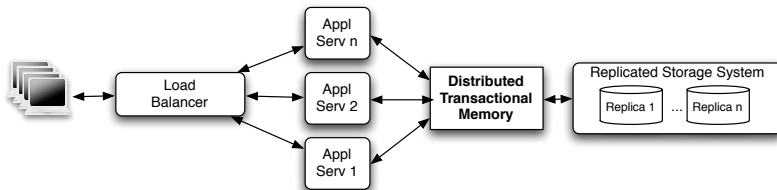
DSTM vs Replicated Databases

- *Similar*: Atomic Broadcast can be used to achieve a global serialization order
- *Different*: The relative overhead of the Atomic Broadcast is bigger

Example Application

Distributed transactional cache for multi-tier applications

- Allows local processing of requests
- Detects both local and remote conflicts
- Alleviates pressure on back-end persistent storage



FenixEDU

University campus management system

- Used in an engineering school in Portugal
- Real system with real scalability and reliability issues

The image displays two overlapping screenshots of the FenixEDU system. The background screenshot shows the main portal of the Instituto Superior Técnico (IST), featuring a navigation menu with 'SOBRE O IST', 'ENSINO', and 'INVESTIGAÇÃO'. Below the menu, there are sections for 'Licenciaturas' and 'CURSO', listing various engineering disciplines and their models. A central banner for 'III Temporada de MÚSICA do Instituto Superior Técnico' is visible. The foreground screenshot shows the 'istid' login page, which includes a 'Login' button, a 'PT | EN' language selector, and a list of authentication methods: 'IST ID', 'Citizen Card', 'Kerberos', and 'European e-ID'. The 'IST ID' method is selected, and the login form contains fields for 'IST ID' (with the value 'ist90476'), 'Password', and a 'Login' button. Below the form, there are links for 'Campus Appstore', 'Cartão Registo', 'Iniciativas + Recursos', and 'Observatório de Empregabilidade Recrutamento'.

Goals

Fault tolerance

Using replication schemes, already studied in other transactional systems (Databases)

Scaling up

Scale up in the number of STM instances to increase performance

Key Idea

- Use optimistic message deliveries to estimate the final transaction certification order
- Expose fresh (although possibly erroneous) data to new transactions
- Reduce the abort rate and detect conflicts earlier

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Distributed STMs

- Distributed Multi Versioning (Manassiev et al.)
- DiSTM (Kotselidis et al.)
- Cluster-STM (Bocchino et al.)

Fault tolerance is not the focus of previous work

Replicated STMs/DBMS

- Active replication without speculation:
 - (Kemme et al.) – uses optimistic total order to speedup commit but does not make speculative results visible
- Active replication with speculation
 - AGGRO (Palmieri et al.) – good for light weight transactions, as all nodes have to execute all transactions
- Certification without speculation
 - D²STM (Couceiro et al.) and ALC (Carvalho et al.)

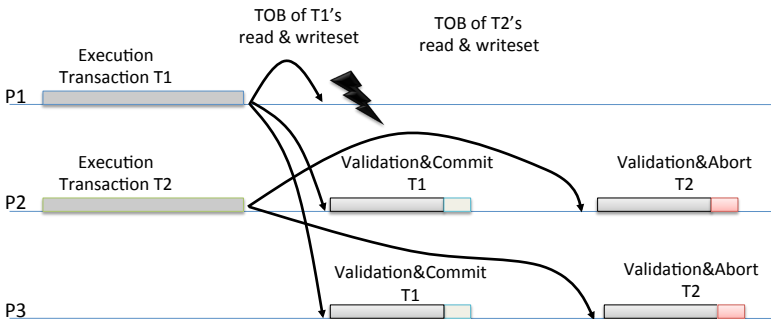
Related Work

	Active Replication	Certification
Non-Speculative	(Kemme et al.)	D ² STM and ALC
Speculative	AGGRO	SCert

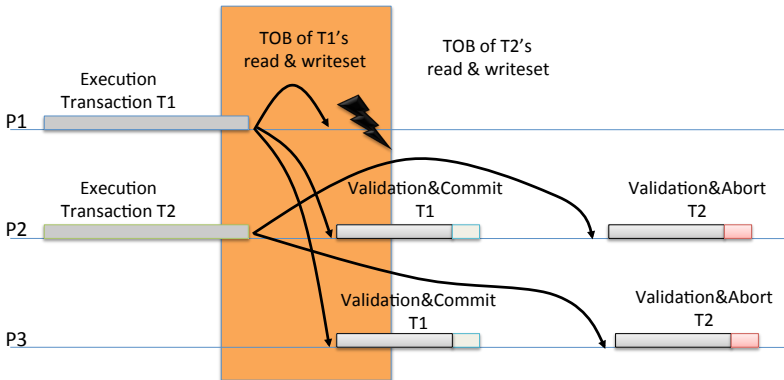
Replication Protocol Based on Certification

- Executes transactions in a single machine optimistically
- Transactions are certified only at commit time
- Exploits *Atomic Broadcast* to ensure replica consistency

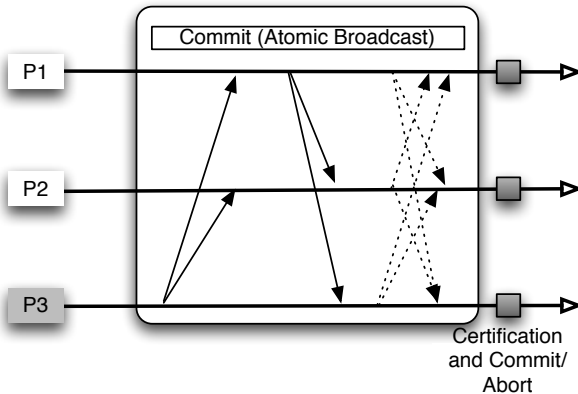
Baseline Replication Protocol



Baseline Replication Protocol



Certification Based Protocol

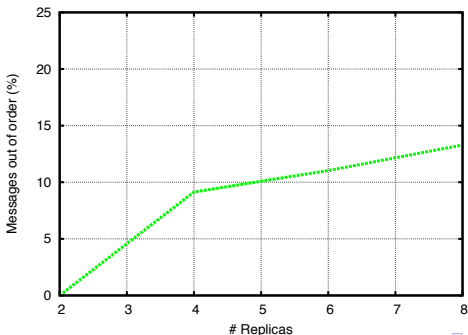


Problems of this Approach

- Loss of efficiency in high conflict scenarios
- Uses a heavy communication procedure (Atomic Broadcast)

Optimistic Atomic Broadcast (OAB)

- Delivers the message twice: an early estimate of the final order and the final order itself
- The estimated order matches the final order with high probability, on LANs



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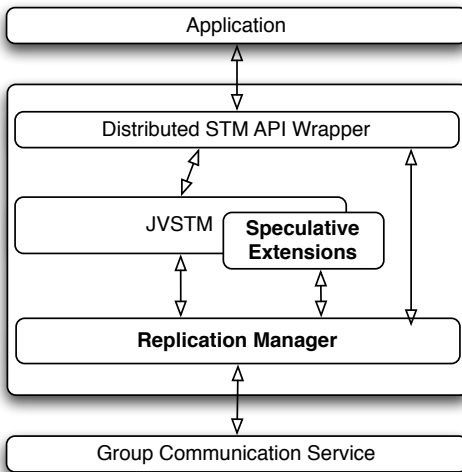
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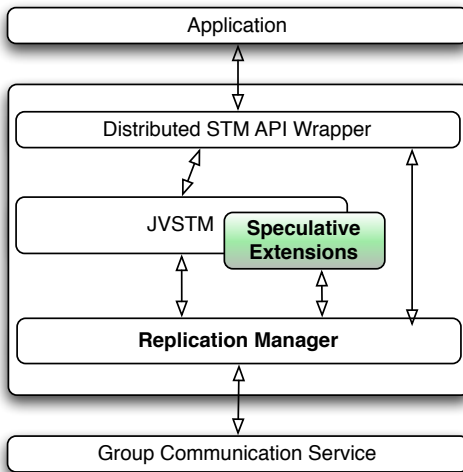
Speculative Certification (SCert)

- Certification based replication protocol
- Exploits *Optimistic deliveries of OAB* to generate fresh (but possibly erroneous) data
- New transactions read the optimistic data snapshots:
 - **Provide** executing transactions with **fresher snapshots**, reducing the probability of aborts
 - **Detect conflicts earlier**, reducing the amount of wasted computation and waiting time

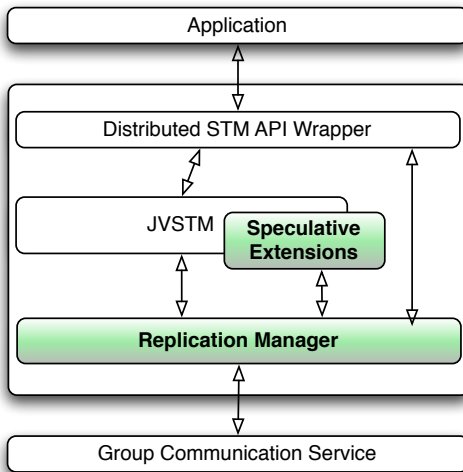
SCert: Architecture



SCert: Architecture



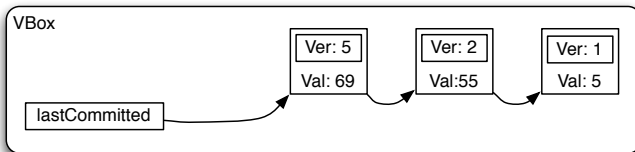
SCert: Architecture



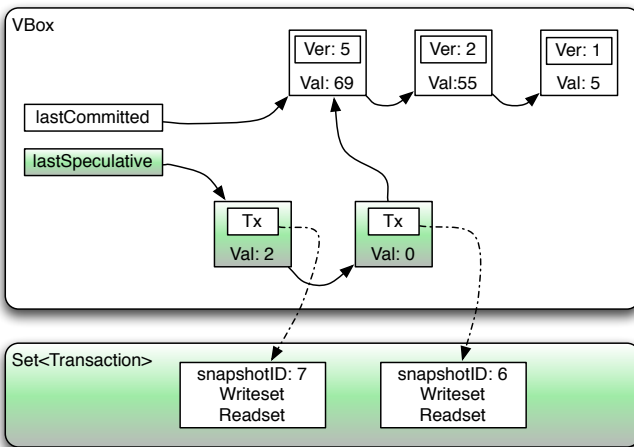
Speculative Extensions

- Provide the appropriate tools to expose speculative committed memory snapshots
- Speculative versions must be maintained
- The API must support speculative commits

JVSTM: Regular VBox



JVSTM Extensions for Speculative Transactions



SCert Replication Protocol (I)

- Transaction executes locally
- Upon Commit, the thread (locally) certifies the transaction and sends OAB
- Upon Optimistic Delivery, the transaction is certified and optimistically committed

SCert Replication Protocol (II)

- Upon Begin of new transactions, the new threads read the most fresh data (committed optimistically or finally)
- Upon Final Delivery:
 - Order matches: the transaction is marked as committed and the thread is unblocked
 - Order does not match: the optimistically committed snapshot is discarded and pending transactions must be re-certified

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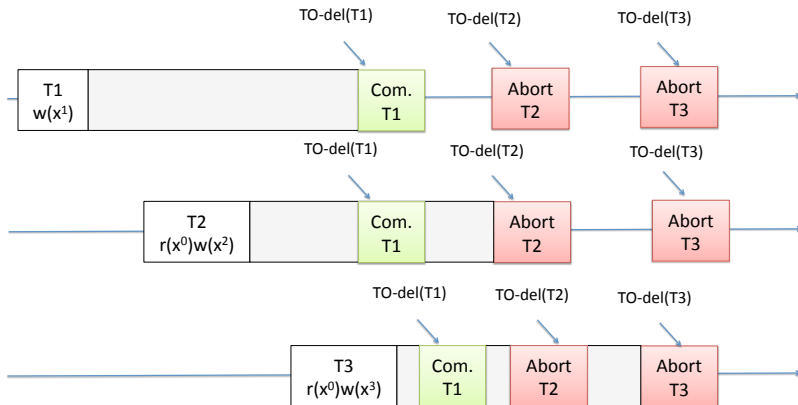
SCert

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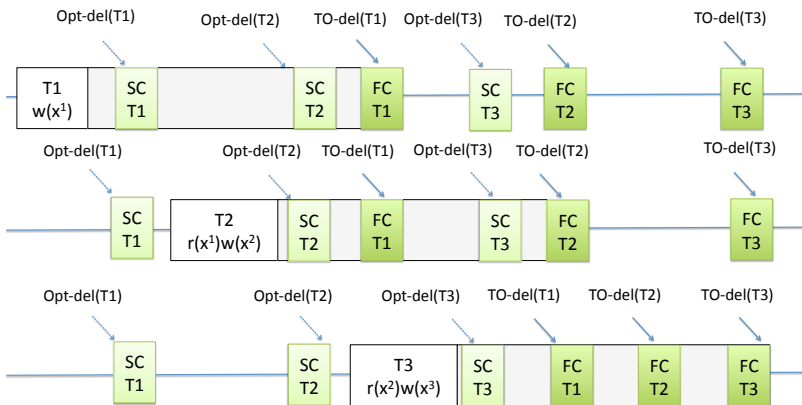
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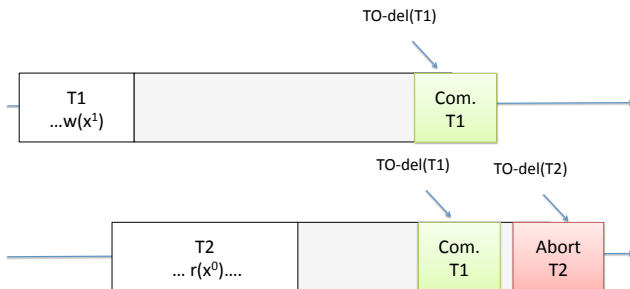
Regular Certification: Cascading Aborts Due to Conflicts



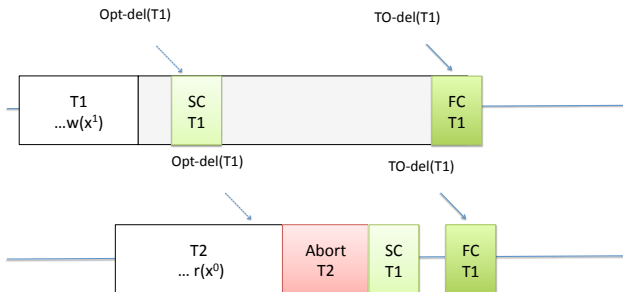
SCert: Cascading Commits



Regular Certification: Wasted Time



SCert: Early Notification



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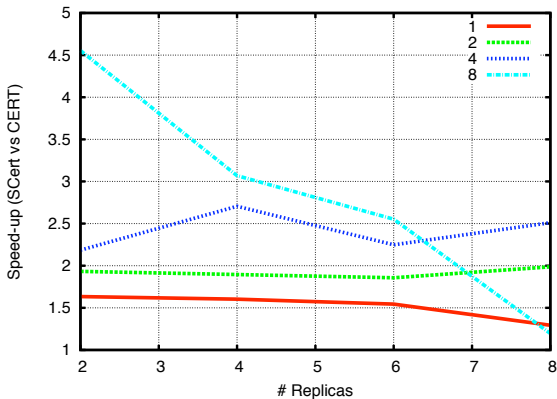
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Bank Benchmark: Full Conflict Scenario

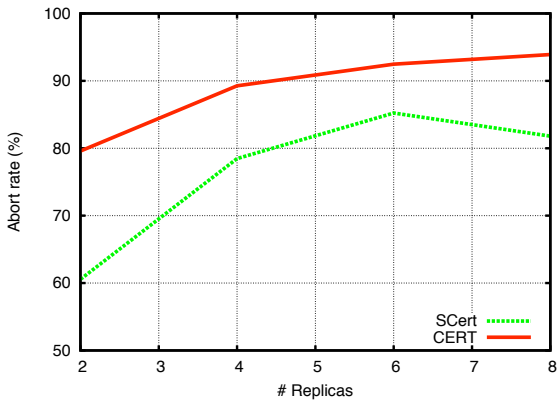
- Goal: worst case
- Replicas accessing the same memory region

Bank Benchmark: Throughput in Worst Scenario



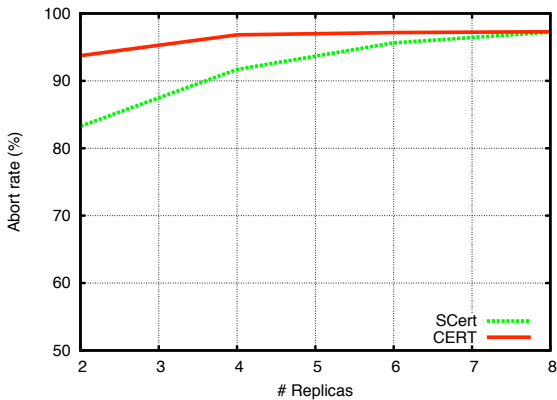
On avg. **1.5x speedup** with one thread and up to **4.5x speedup** with 8 threads per replica

Bank Benchmark: Abort Rate



1 thread per replica

Bank Benchmark: Abort Rate

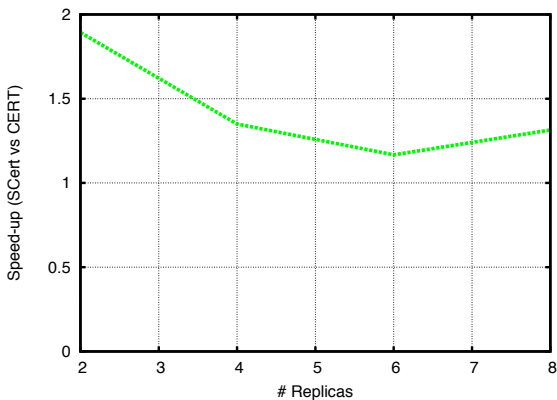


8 threads per replica

STMBench7 Benchmark: Scenario

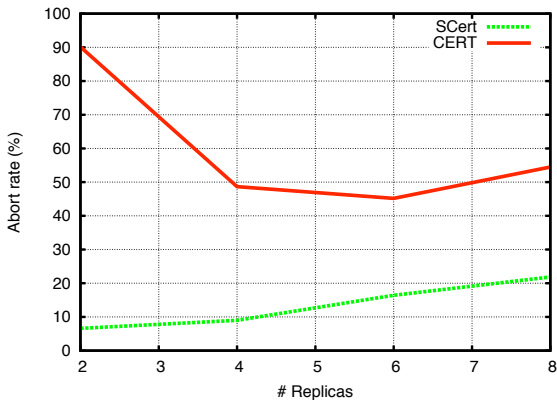
- Goal: more complex benchmark
- Richer benchmark featuring a number of operations with different levels of complexity over an object-graph with millions of objects
- Number of machines between 2 and 8
- Number of threads fixed to 2

STMBench7 Benchmark: Speedup



Almost **twice speedup** with a low number of replicas

STMBench7 Benchmark: Abort Rate



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- Reduce the number of transactions that read stale data
- Allows early detection of conflicts among transactions
- Performance improvements are achieved by exploiting optimistic deliveries of OAB
 - Up to 4.5x speed-ups

Thank you!

Questions?